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**A LOGIC TABLE ABSTRACTION LAYER FOR  
ACCESSING CONFIGURATION INFORMATION**

5

**Technical Field**

The invention relates generally to object-oriented data processing and system management, and more particularly to a logic table abstraction layer used to access configuration information in a catalog environment.

**Background of the Invention**

10 In a distributed computer environment, client computer systems and server computer systems can share data and executable program code, including applications, services, and resources. In order to facilitate the sharing of these various elements, attribute-based programming allows a developer to specify the services and resources required by an application by setting properties (or “attributes”) of each application or component rather than implementing or calling those services directly from the implementation code. Attributes include a particular  
15 set of configuration information that is made available to various callers in an attribute-based programming environment.

20 Configuration information may be stored on the various machines in the distributed network and relates to specific properties of applications, components, services, and other resources available to that machine. In existing approaches, a system “registry” has been used to store configuration information for a particular machine. However, in existing approaches, a programmer is required to access and manipulate registry information directly, introducing undesirable program complexity and exposing the registry to corruption by improper

programming. Moreover, the distribution of configuration information among multiple datastores (i.e., in addition to the registry) and data formats is not accommodated by existing approaches, particularly if the location and format of data is expected to evolve over time. In current approaches, the implementation code itself must be altered in order to handle location and format changes to configuration information. Accordingly, existing approaches lack location and format independence that can provide the desired flexibility for storage and access to configuration information in a computer system.

### **Summary of the Invention**

In accordance with the present invention, the above and other problems are solved by providing a logic table object that abstracts access to underlying configuration data. A logic table object can provide supplemental logic and trigger external operations not provided by underlying table objects and datastores. Furthermore, a logic table object can synthesize data not available from underlying table objects. In addition, a logic table object maps coordinates of a logic level table to one or more coordinates of underlying tables.

A computer program product providing domain-specific access to configuration information sourced by at least one datastore, wherein the domain-specific access is substantially specified by at least one input parameter, is provided. A first level table object is instantiated in accordance with a input parameter. The first level table object includes a first table-oriented interface having a first table-oriented method. A call is received to the first table-oriented method in the first level table object. A logic component module executes, responsive to the call, if the first level table object provides domain-specific logic corresponding to the first table-oriented method. The call is delegated to a corresponding table-oriented method of a lower-level

table object to which the first level table object is bound, if the first level table object depends on the lower-level table object to completely service the call.

A logic table object, executable by a computer, provides domain-specific access to configuration information sourced by at least one datastore, wherein the domain-specific access is substantially specified by at least one input parameter. A table-oriented interface includes a table-oriented method accessible by a caller to access the configuration information and receives a call from the caller to the table-oriented method. A logic component module providing domain-specific logic to the table-oriented method. An interception/delegation module executes the domain-specific logic of the logic component module, responsive to receipt of the call, and further delegates the call to a corresponding table-oriented method of a lower-level table object to which the logic table object is bound, if the logic table object depends on the lower-level table object to completely service the call.

### **Brief Description of the Drawings**

FIG. 1 illustrates a pictorial representation of a suitable client-server computing environment in which an embodiment of the present invention may be implemented in both clients and servers.

FIG. 2 depicts an exemplary client/server architecture employing COM+ catalogs (Component Object Model) in accordance with the present invention.

FIG. 3 illustrates an exemplary system for implementing the invention in an embodiment of the present invention.

FIG. 4 depicts various examples of table system configurations in an embodiment of the present invention.

FIG. 5 depicts a logic table object in an embodiment of the present invention.

FIG. 6 depicts a logic table object including a cache and individual logic component modules in an embodiment of the present invention.

FIG. 7 illustrates exemplary results of data consolidation, triggering, and data synthesis in an embodiment of the present invention.

FIG. 8 illustrates a flowchart of operations for providing supplemental table logic in an embodiment of the present invention.

FIG. 9 illustrates a flowchart of operations for mapping table data an embodiment of the present invention.

FIG. 10 illustrates a flowchart of operations for synthesizing data in an embodiment of the present invention.

### **Detailed Description of the Invention**

An embodiment of the present invention includes a logic table object for accessing configuration information sourced by one or more datastores. In a preferred embodiment, logic table objects are employed in a catalog environment of an attribute-based programming model; however, embodiments of the present invention can be broadly applied to access any type of data. A logic table object may provide supplemental logic, consolidate multiple underlying tables, map between disparate table coordinates, trigger external operations, and synthesize data for inclusion in a virtual table presented to a caller.

FIG. 1 is a pictorial representation of a suitable client-server computing environment in which an embodiment of the present invention may be implemented in both clients and servers. In a computing network 100, client computer systems 102, 104, 106 and 108 are connected to

server computer systems 110, 112, 114 and 116 by a network connection 118. Additionally, client computer 120 is connected to server computer 110 via a communication link, such as the Internet 122 or a local area network. Since the server 110 is connected via the network connection 118 to the other servers 112, 114 and 116, the client computer 120 is also connected and may access information on the other servers 112, 114 and 116, and clients 102, 104, 106, and 108, as well as other computer systems coupled to the network 100.

The client computer systems 102, 104, 106, 108 and 120 operate using at least some of the information and processes available on at least one of the servers 110, 112, 114 and 116 as well as other computer systems coupled to the network 100. Each client is preferably a complete, stand-alone computer and offers the user a full range of power and features for running applications. The clients 102, 104, 106 and 108, however, may be quite different from the other clients as long as they can communicate via the common interface 118.

The servers 110, 112, 114 and 116 are preferably computers, minicomputers, or mainframes that provide traditional strengths offered by minicomputers and mainframes in a time-sharing environment (e.g., data management, information sharing between clients, and sophisticated network administration and security features). The client and server machines work together to accomplish the processing of the executed application. Working together in this manner increases the processing power and efficiency relating to each independent computer system shown in FIG. 1.

Typically, a client portion or process of an application executed in the distributed network 100 is optimized for user interaction whereas a server portion or process provides the centralized, multi-user functionality. However, each client computer 102, 104, 106, 108 and 120 can perform functions for other computers, including the clients and servers, thus acting as a “server” for

those other computer systems. Similarly, each of the servers 110, 112, 114 and 116 can perform functions and relay information to the other servers, such that each server may act as a “client” requesting information or services from another computer in particular circumstances.

Therefore, the term “client,” as used hereinafter refers to any computer system making a call or request of another computer system and the term “server” is the computer system servicing the request.

As part of the sophisticated network administration, each computer is able to access configuration information related to applications and resources available on the other computers in the network 100. The configuration information is located within memory or persistent storage on each computer system, i.e., in a datastore. Additionally, each computer system may have more than one datastore of configuration information that must be accessed by the other computer systems. Moreover, the different datastores may each have different data types or formats. In order to access configuration information from these many and various computer datastores, a client, i.e., the system or process making the request for information, communicates with a “catalog” interface on the computer system.

FIG. 2 depicts an exemplary client/server architecture employing COM+ catalogs in accordance with the present invention (COM is an acronym for Component Object Model). A COM+ Catalog is a virtualized database of COM+ applications and their services, with runtime and configuration-time abstraction layers for using and manipulating configuration information. An embodiment of the present invention, for example, may be employed in a component-based programming model of a transaction processing runtime environment for developing, deploying, and managing high-performance, scaleable, and robust enterprise Internet and intranet server applications.

A “component” is software containing classes that may be created and exposed as “objects” (i.e., self-contained programmed entities that consist of both data and functions to manipulate the data) for use by another application. A component can also use objects exposed by another application. For example, a developer can create an application using ActiveX components that can be updated and managed easily as in-process DLLs (Dynamic Link Libraries). The DLLs are then installed into the COM environment for execution within the application. Components can be developed specifically for a developer’s single application, developed for use with multiple applications, or purchased from a third party.

COM technology allows a piece of software to offer services to another piece of software by making those services available as “COM objects”. COM is a foundation for an object-based system that focuses on reuse of interfaces. It is also an interface specification from which any number of interfaces can be built. Each COM object is an instance of a particular class and supports a number of interfaces, generally two or more. Each interface includes one or more methods, which are functions that can be called by the objects’ clients.

COM+ technology is an extension of COM technology that includes a new runtime library that provides a wide range of new services, such as dynamic load balancing, queued components, an in-memory database, and events. COM+ technology maintains the basics of COM technology, and existing COM-based applications can continue to work unchanged in a COM+ environment.

An object implemented to comply with COM+ is referred to as a “COM+ object”. A component that includes one or more classes that may be instantiated as a COM+ object is referred to as a “COM+ component”. Each COM+ component has attributes, which can be set in a component (or type) library. Attributes are a form of configuration data required by many



software components to execute correctly and completely. An application that includes COM+ components is referred to as a "COM+ application". When a component is made part of a COM+ application, its component (or type) library is written into a COM+ catalog. When an object is instantiated from that component, the attributes in the COM+ catalog are examined to determine the object context that contains properties for the object. Based on the object context, other services required by the object are provided. In this manner, a developer can merely identify in the attributes the additional functionality required by the object, and based on the object's attributes, the appropriate other services that are available within the system, or the accessible network, are executed to provide that functionality.

In FIG. 2, a client computer 200 is coupled via a network to one or more remote computers (e.g., a computer 202 and a server 204). Although the embodiments of the present invention are illustrated and described herein relative to multiple computer systems coupled by a computer network or other communications connection, it is to be understood that an embodiment of the present invention may be employed in a stand-alone computer system to provide access to configuration information in the system.

A client application 206 executes on the client computer 200 to access a server application 208 executing on the server 204. For example, the server application 208 may include a database application that receives a query from the client application 206 and accesses a customer database (not shown) for all customer data satisfying the query. During operation, the server application 208 may require configuration data recorded in a datastore (such as datastores 214 or 216). For example, a transaction server application can determine the security level of a user according to a "role" assigned to the user by an administrator or other means. Accordingly, the transaction server application might query a role definitions database to validate

the user's access to a transaction database (not shown). In another example, the server application 208 accesses configuration information to verify that required services are available for its execution.

To obtain configuration information in the illustrated embodiment, the server application 208 accesses a runtime catalog 210 running on the server 204. The runtime catalog 210 causes one or more table object dispensers to create catalog table objects (shown generally as table system 218) providing the required configuration data in a table to the server application 208. A "table object" includes an object that provides a caller with access to underlying data, presenting that data in virtual "table" format through a defined table interface. A table object may also provide its own functionality, read and write caching and the triggering of external events, in addition to other features. The table data is accessed by a caller (e.g., a catalog server, a runtime catalog, or an overlaying logic table object) by way of a table-oriented interface, preferably including table cursor methods. In the exemplary embodiment, the runtime catalog 210 accesses configuration data in the datastores 214 and 216 through layers of abstraction provided by the table system 218 (i.e., including logic table objects (LT), such as logic table object 220, and data table objects (DTs), such as data table object 222).

A globally unique database ID (identifier) called a "DID" identifies each catalog database. A given DID guarantees a minimum well-defined set of catalog tables, each table being identified by and complying to the rules of a table ID (TID). A DID is a datastore-independent identity, meaning that the tables of that database can be distributed among multiple datastores. Examples of datastores include the registry, type libraries, SQL (structured query language) Servers, and the NT Directory Service (NT DS), whereas examples of databases include: server group databases, download databases, and deployment databases.

A data table object, such as data table object 222, is a datastore-dependent table object that exposes a table cursor into a particular datastore. The table cursor provides a well-defined table-oriented interface into the datastore while hiding the location and format of the underlying datastore itself. For example, a caller can use a table cursor to navigate through the rows of a column in a table presented to the caller by a table object.

Each data table object is bound to a particular datastore accessible within the computer. For example, a data table object may be bound to the registry to provide the registry data in table form to a higher level (e.g., an overlaid logic table object, catalog server object, or runtime catalog). Another data table object may be bound to the NT Directory Services to provide directory configuration data to a higher level. As shown by data table objects 238 and 240, multiple data table objects may be created for a single datastore (e.g., data table objects 238 and 240 are created by different logic tables objects to provide access to the same datastore 242).

The data table object 222 populates one or more internal caches with read or write data associated with the datastore 214. Queries to the datastore 214 are serviced by the cache or caches through the data table object's table interface. Using at least one "update" method, data in the read cache of data table object 222 may be refreshed from the datastore 214 and data in a write cache may be flushed to the datastore 214. Data table objects are described in more detail in U.S. Patent Application No. \_\_\_\_\_, entitled "DATA TABLE OBJECT INTERFACE FOR DATASTORE," assigned to the assignee of the present application, filed concurrently herewith and incorporated herein by reference for all that it discloses and teaches.

A logic table object, such as logic table object 220, presents domain-specific table data by logically merging or consolidating table data from multiple data table and/or logic table objects, supplementing table functionality, and/or synthesizing data into the table, in accordance with a

given type of configuration information requested (e.g., configuration information for Components, Applications, etc.). The domain-specific nature of the table data is preferably defined by at least one input parameter, including without limitation a database ID, a table ID, a query parameter, or a level of server parameter). Logic table objects in a COM+ Catalog environment are type-independent abstraction layers between a caller (such as the runtime catalog 210) and one or more datastores (such as datastores 214 and 216) containing configuration information. A logic table object typically sits atop one or more data table objects and introduces domain-specific rules and processes to the underlying data table objects, although other configurations of table systems are possible (see FIG. 4).

More specifically, a logic table object can logically merge or consolidate configuration data from multiple data table and/or logic table objects into a single table based on predetermined logic (e.g., according to type). Furthermore, a logic table object can supplement data table object functionality by intercepting interface calls from a client and adding to or overriding the underlying table object functionality (e.g., adding validation or security). Additionally, a logic table object can synthesize data that is not available from the underlying datastores or tables and present the synthesized data as part of the table.

The foregoing discussion has described the COM+ Catalog environment as used at runtime by an application. An alternate use of a COM+ Catalog occurs at configuration-time and may employ one or more catalog server objects (CS) and one or more client tables. During configuration, an administration tool, such as Microsoft's Component Services administration tool or COMAdmin Library, is used to create and configure COM+ applications, install and export existing COM+ applications, manage installed COM+ applications, and manage and configure services locally or remotely. Accordingly, in addition to the illustrated embodiments,

an embodiment of the present invention may be employed by a local administration tool managing an application running on a remote computer system.

The exemplary administration tool 224 executes on the client computer 200 in FIG. 2. An alternative administration tool (such as administration tool 250) can execute on another computer (such as server 204) to configure applications and services executing in the computer. A catalog server object, such as catalog server objects 226, 228, and 230, manages configuration information on its computer. All administration requests to a computer, whether local or from another computer, go to a catalog server object on that computer, preferably through one or more abstraction layers, including client table objects and logic table objects.

A client table object (CT) is analogous to a data table object that binds to a particular local or remote catalog server object instead of a datastore, presenting the configuration information marshaled by a catalog server object in table form to the caller, such as the administration tool 224. The local catalog server object 226 manages configuration data locally on the client computer 200, while the remote catalog server object 228 service catalog requests from the client table object 232 for configuration information on its remote computer.

“Remote” does not necessarily imply that a remote computer geographically distant from a local computer. Instead, remote merely indicates a cross-computer boundary, which may be bridged by a data bus, a network connection, or other connection means.

To access available catalog data in the illustrated exemplary embodiment, the administration tool 224 optionally causes a logic table object 234 to be created, which in turn causes client table objects 232 and 236 to be created for accessing available catalog server objects 226, and 228. The local catalog server object 226 and the remote catalog server object 228 marshal the configuration information stored within their corresponding computers by

causing creation of underlying table systems and transferring the data back to the client table objects 232 and 236 for presentation as table data to the logic table object 234, which logically merges the configuration information and presents the configuration information to the administration tool 224 in table format. In the illustrated embodiment, multiple domain-specific logic table objects (such as logic table object 234) can reside between the client table objects 232 and 236, and the administration tool 224. Alternatively, the administration tool 224 may cause only a single client table object (with or without overlaying logic table objects) to be created to access a single catalog server object on a local or remote computer.

With reference to FIG. 3, an exemplary computing system for embodiments of the invention includes a general purpose computing device in the form of a conventional computer system 300, including a processor unit 302, a system memory 304, and a system bus 306 that couples various system components including the system memory 304 to the processor unit 300. The system bus 306 may be any of several types of bus structures including a memory bus or memory controller, a peripheral bus and a local bus using any of a variety of bus architectures. The system memory includes read only memory (ROM) 308 and random access memory (RAM) 310. A basic input/output system 312 (BIOS), which contains basic routines that help transfer information between elements within the computer system 300, is stored in ROM 308.

The computer system 300 further includes a hard disk drive 312 for reading from and writing to a hard disk, a magnetic disk drive 314 for reading from or writing to a removable magnetic disk 316, and an optical disk drive 318 for reading from or writing to a removable optical disk 319 such as a CD ROM, DVD, or other optical media. The hard disk drive 312, magnetic disk drive 314, and optical disk drive 318 are connected to the system bus 306 by a hard disk drive interface 320, a magnetic disk drive interface 322, and an optical drive

interface 324, respectively. The drives and their associated computer-readable media provide nonvolatile storage of computer readable instructions, data structures, programs, and other data for the computer system 300.

Although the exemplary environment described herein employs a hard disk, a removable magnetic disk 316, and a removable optical disk 319, other types of computer-readable media capable of storing data can be used in the exemplary system. Examples of these other types of computer-readable mediums that can be used in the exemplary operating environment include magnetic cassettes, flash memory cards, digital video disks, Bernoulli cartridges, random access memories (RAMs), and read only memories (ROMs).

A number of program modules may be stored on the hard disk, magnetic disk 316, optical disk 319, ROM 308 or RAM 310, including an operating system 326, one or more application programs 328, other program modules 330, and program data 332. A user may enter commands and information into the computer system 300 through input devices such as a keyboard 334 and mouse 336 or other pointing device. Examples of other input devices may include a microphone, joystick, game pad, satellite dish, and scanner. These and other input devices are often connected to the processing unit 302 through a serial port interface 340 that is coupled to the system bus 306. Nevertheless, these input devices also may be connected by other interfaces, such as a parallel port, game port, or a universal serial bus (USB). A monitor 342 or other type of display device is also connected to the system bus 306 via an interface, such as a video adapter 344. In addition to the monitor 342, computer systems typically include other peripheral output devices (not shown), such as speakers and printers.

The computer system 300 may operate in a networked environment using logical connections to one or more remote computers, such as a remote computer 346. The remote

computer 346 may be a computer system, a server, a router, a network PC, a peer device or other common network node, and typically includes many or all of the elements described above relative to the computer system 300. The network connections include a local area network (LAN) 348 and a wide area network (WAN) 350. Such networking environments are  
5 commonplace in offices, enterprise-wide computer networks, intranets, and the Internet.

When used in a LAN networking environment, the computer system 300 is connected to the local network 348 through a network interface or adapter 352. When used in a WAN networking environment, the computer system 300 typically includes a modem 354 or other means for establishing communications over the wide area network 350, such as the Internet. The modem 354, which may be internal or external, is connected to the system bus 306 via the  
10 serial port interface 340. In a networked environment, program modules depicted relative to the computer system 300, or portions thereof, may be stored in the remote memory storage device. It will be appreciated that the network connections shown are exemplary, and other means of establishing a communication link between the computers may be used.

In an embodiment of the present invention, the computer system 300 stores the configuration data and implementation code providing the catalog infrastructure and disclosed and claimed herein in accordance with the present invention. The catalog infrastructure has without limitation one or more datastores, catalog servers, runtime catalogs, server applications, administration tools, dispensers, and wiring databases. Specifically, one or more dispensers,  
15 preferably including a table dispenser and a table object dispenser, provide a table object to a caller providing location and type independent access to configuration information stored in one or more datastores.



Preferably, table objects for accessing one or more datastores in a computer system are obtained via one or more table dispensers or table object dispensers. To access one or more datastores, a caller obtains a table object by passing input parameters to a table dispenser. The table dispenser references a wiring database to determine an appropriate configuration of table objects needed to return the desired table object to the caller. Dispensers are described in more detail in U.S. Patent Application No. \_\_\_\_\_, entitled "OBTAINING TABLE OBJECTS USING TABLE DISPENSERS", filed concurrently herewith and incorporated herein by reference for all that it discloses and teaches.

FIG. 4 depicts various examples of table systems in embodiments of the present invention. Logic table and data table objects are described in the description of FIG. 2 and the incorporated references. With regard to a table system 400, a caller 410 (as well as other callers in FIG. 4) may be a catalog server, a runtime catalog, or another object requiring abstracted access to a datastore. To initiate access to requested information, the caller 410 provides input parameters, such as a database ID, a table ID, query parameters, and a level of service parameter, relating to the configuration information it is requesting. A table dispenser (see the table dispenser 502, for example, in FIG. 5A) returns to the caller 410 a pointer to a table object, in this case a single data table object 412 bound to a datastore 414. Through a table interface accessible via the pointer to the data table object 412, the caller 410 can access tabularized configuration data (i.e., a data level table) originating from the datastore 414.

With regard to a table system 402, the table dispenser provides a caller 416 with an interface of a logic table object 418, which overlays a data table object 420. The data table object 420 is bound to a datastore 422 and provides access to a data level table of configuration data originating from the datastore 422 to the logic table object 418. Through a table interface

provided to caller 416, the logic table object 418 can present to the caller a logic level table of configuration information, including without limitation (1) a remapping (i.e., an alternate table configuration) of the data provided by data table object 420; (2) supplemental functionality (e.g., validation of data); and (3) synthesized data (e.g., data not resident in datastore 422, but instead, data derived or calculated from data in datastore 422 or another source). The logic table object 418 can also trigger external operations.

With regard to a table system 404, two levels of logic table objects (i.e., logic table objects 426 and 428) are positioned between a caller 424 and a data table object 430, which is bound to a datastore 432. Preferably, functionality is modularized using multiple logic table objects. For example, the logic table object 426 may be responsible for enforcing security constraints on accesses to configuration data, whereas the logic table object 428 may validate data before writing configuration data to the datastore 432. Other functional combinations are possible at the discretion of the developer. In an embodiment of the present invention, the combinations of logic table and data table objects required to satisfy a requested database ID and table ID are specified in a wiring database accessed by the table dispenser.

With regard to a table system 406, a caller 434 has access to configuration data through a logic table object 436 without an underlying data table object or datastore. The logic table object 436 may provide table-based synthesized data to the caller or otherwise provide or trigger functionality outside the scope of the catalog's tables. For example, the logic table object 436 may intercept calls to an unsupported datastore and return errors to the caller 434. Alternatively, the logic table object 436 may translate or remap table data originally provided by the caller 434 or an external source, rather than by a datastore.

With regard to a table system combination 408, a caller 438 gains access to configuration data originally stored in or derived from datastores 450, 452, or 454. A logic table object 440 logically merges or consolidates data from a logic table object 442, data table object 444, which is bound to datastore 452, and data table object 446, which is bound to datastore 454. The logic table object 442 overlays a data table object 448, which is bound to datastore 450. In this configuration, the logic table object 440 logically merges data from the underlying catalog tables and presents the configuration data as a logic level table to the caller 438.

FIG. 5 depicts a logic table object in an embodiment of the present invention. A logic table object 500 presents a table interface 502 to a caller 503. The table interface 502 is compatible with the interface presented by the data table objects 504. The table interface 502 is compatible with, and preferably identical to, the data table interfaces 514. The table interface 502 includes methods for navigating rows of a table, retrieving configuration information values from columns of a logic level table 506, reading metadata from the logic level table or a column thereof, deleting rows from the logic level table 506, inserting and updating rows from the logic level table 506, updating a datastore, populating a read cache from a datastore, and other advanced operations. In an embodiment of the present invention, metadata is read to define a schema of a logic level table or data level table.

Cursor methods are a type of table-oriented method supported by a logic table embodiment of the present invention. Table 1 includes descriptions of cursor methods in an exemplary table interface:

Method	Description
PopulateCache	Populates the read cache from a datastore or an underlying table object, using the database, the table, and the query

	specified to the table dispenser. The previous cache contents, including pending changes, are discarded. The row cursor is restarted. The table dispenser makes the first (and typically the only) call to populate the read cache before providing the table cursor to its caller.
UpdateStore	Writes all pending changes in the write cache to the datastore and then clears the write cache. The UpdateStore method is not implemented in a read-only cache for which it preferably returns an error.
UpdateReadCache	Updates the read cache with the pending changes in the write cache. The row cursor is restarted. The contents of the datastore and the write cache remain unchanged. The UpdateReadCache method is not implemented on a read-only cache for which it preferably returns an error.
RestartRow	Restarts the row cursor just prior to the first row in the read cache.
GetNextRow	Moves the row cursor from the current row to the next row in the read cache. When preceded by RestartRow, the GetNextRow method moves the cursor to the first row in the read cache. When the cursor is on the last row or a new row, the GetNextRow method does not move the cursor and preferably returns an error.
MoveToRowByIdentity	Moves the row cursor to the unique row in the cache matching the identity specified in the method call. If the row does not exist, the MoveToRowByIdentity method returns an error and does not move the cursor.
MoveToNewRow	Moves the cursor to a new row held by the cursor. The MoveToNewRow method is used in combination with the SetRow method, which inserts the row into the write cache. All non-default columns are first set with SetColumn. The SetRow method must be called before the cursor moves. The MoveToNewRow returns an error on a read-only cache.
DeleteRow	Marks a current row deleted in the write cache. The DeleteRow method returns an error on a read-only cache.
SetRow	Updates the write cache (not the datastore) with the changes made to the current row. If the row is from MoveToNextRow, but the row is already identified in a read cache, the write cache is still updated with changes to the row. The SetRow method returns an error on a read-only cache.
SetColumn	Prepares to set column i of the current row using the data supplied in the input parameters of the method call. The change is held by the cursor until SetRow is called. If the data type of the supplied data only requires four bytes, the caller passes the data directly; otherwise, the caller passes the

	data by reference and, therefore, must retain the data until SetRow is called. If column i is out of range, the SetColumn method returns an error. The SetColumn method also returns an error on a read-only cache.
GetColumn	Gets the data for column i of the current row. If column i is out of range, then GetColumn returns an error.
CloneCursor	Supplies another cursor, initially at the same location in the read cache as the current cursor.

Table 1 - An Exemplary Table Interface Presented by a Table Object

In order to implement the table interface 502, a logic table object preferably intercepts all cursor method calls from the caller 503. "Intercepting" involves receiving a cursor method call to a table object and performing logical operations independent of operations that could be obtained by merely delegating the cursor method call to an underlying table object. In some cases, intercepting can completely replace delegation to the corresponding cursor method in an underlying table object. More specifically, the caller calls a given cursor method indicated by a pointer in a virtual table (vtable) associated with the logic table object 500, passing appropriate parameters to the cursor method. A virtual table containing the addresses (pointers) for the methods and properties of each table object.

The operations that occur on the logic table object side of the table interface 502 are transparent to the caller 503. As such, in response to each call, the logic table object 500 can supplement the functionality of underlying table objects, directly delegate processing to one or more underlying tables, or entirely provide its own functionality to complete the call.

In an embodiment of the present invention, a logic table object 500 provides an abstraction layer between the caller 503, which is presented with a logic level table 506, and one or more data table objects 504, which are bound to respective datastores 510. The logic table object 500 includes an interception/delegation module 508, which associates cursor methods

and/or data in the logic table object 500 to the logic level table 506. The interception-delegation module 508 of the illustrated embodiment may also delegate calls to the table interfaces of one or more underlying data table objects 504. In alternative embodiments, one or more other logic table objects can be underlying table objects. By delegating, the logic table object 500 effectively passes the call through to one or more underlying data table objects (e.g., by calling a corresponding cursor method of the underlying table object with equivalent or otherwise appropriate calling parameters). By intercepting, however, the logic table object 500 can provide additional logic (e.g. supplemental functions, synthesized data, mapping, or consolidation) responsive to the call.

While it is possible to delegate by passing a call to an underlying cursor method, an optimized method of delegation avoids the overhead of the unnecessary “second” call. When direct delegation is desired (i.e., without mapping, supplemental logic, or synthesis), a pointer directly to the cursor method of an underlying table object can replace the pointer to the logic table object’s cursor method in the vtable. In this manner, although the caller anticipates it is calling a cursor method in the logic table object, the pointer in the vtable actually directs processing to the cursor method in the underlying table object. When the cursor method call completes, processing returns directly to the calling process.

Depending on the specific implementation of the logic table object 500, the additional logic triggered by an interception event and provided by one or more logic component modules 512 may include logical merging or consolidation of configuration information from multiple underlying data level tables 505 and/or other underlying logic level tables (not shown). Logic table objects may additionally or alternatively supplement underlying table functionality by intercepting cursor method calls from a client and adding to or overriding the underlying table

object's functionality. Alternately or additionally, the logic table object 500 can synthesize data that is not available from the underlying datastores or table objects and present the synthesized data as part of the logic level table 506 presented to the caller 503.

FIG. 6 depicts a logic table object including a cache and individual logic component modules in an embodiment of the present invention. The cache may include a read cache, a write cache, or both. A logic table object 600 in an embodiment of the present invention includes a cache 602, an interception/delegation module 604, and logic component modules 606. The logic component modules 606 comprise one or more synthesizing modules 608, one or more supplemental logic modules 610, or one or more mapping modules 612 and a mapping lookup table 614. It should be understood that many combinations of the cache 602, the synthesizing module 608, the supplemental logic module 610, and the mapping module 612 are possible within the scope of the present invention. That is, depending on the requirements of the logic table object, a logic table dispenser may include the cache and all of the modules, no cache and one or more modules, some intermediate combination of the cache and modules, or only one of the modules or cache.

Preferably, the mapping lookup table 614 is used in combination with a mapping module 612. However, alternate mapping methods may be used in another embodiment of the present invention including, without limitation of, those implemented by IF-THEN-ELSE constructs or CASE tables.

In a preferred embodiment, data structure elements corresponding to each row and column element (e.g. coordinate) of the logic level table 616 are recorded in a memory structure as a mapping lookup table 614. The data structures in the mapping lookup table 614 preferably include mapping instructions, such as an identifier of the data table object corresponding to

coordinates of the logic level table 616, the corresponding coordinates of the underlying table object to which the logic level coordinates corresponds, and/or a pointer to additional logic.

Accordingly, upon receiving a cursor method call relating to a coordinate in logic level table 616, the interception/delegation module 604 calls a mapping module 612 to determine the mapping to a corresponding coordinate in an underlying data table 618. The mapping module 612 locates the data structure corresponding to the coordinate of the logic level table 616 and returns to the interception/delegation module 604 a pointer and coordinates to the corresponding underlying table object 618, as well as optional additional data associated with the coordinates in the underlying table object. It should be understood that more than one underlying table object and coordinate combination may be returned by the mapping module if the cursor method applies to multiple underlying table objects. The interception/delegation module 604 then calls a cursor method in each corresponding table object using the pointer or pointers returned from the mapping module 612.

With regard to the supplemental logic module 610, the interception/delegation module 604 intercepts a call from the caller and calls supplemental logic module 610 to provide additional logical operations. The supplemental logic can consist of multiple stages, that is, the supplemental logic module 610 can pre-process or post-process a delegation to one or more underlying table objects 618. Alternatively, the supplemental logic module 610 can completely replace a cursor method of an underlying table object 618, foregoing delegation and returning to the caller (without calling a cursor method in an underlying table object).

Examples of supplemental logic include without limitation enforcing complex relationships among column values in a row when a caller attempts to change the values in a column, filtering server-side row or column reads depending on the security level of a caller,



enforcing and managing complex relationships among different tables as those tables change, and triggering external functionality that lies outside the scope of the catalog tables, responsive to predetermined table changes. A complex relationship between column values might require that a component release date in one column not change unless the version (or “build”) designation changes in another column for the same component (i.e., row).

In an embodiment of the present invention, one or more synthesizing modules 608 correspond to one or more coordinates in the logic level table 616 and provide or derive data not available from an underlying table object or datastore. For example, a commonly used catalog table (e.g., corresponding to “components”) includes data logically merged from multiple data level tables and datastores. Each row of a data level table corresponds to a component in the system and its corresponding properties. A synthesizing module 608, however, may be implemented to provide an index for each component having a given property. The caller can then access each index via a coordinate in the logic level table 616.

In an embodiment of the present invention, the call associated with the “index” coordinate of the logic level table 616 is intercepted by the interception/delegation module 604, which calls the mapping module 612 to determine the underlying table object and coordinate. A data structure in the mapping lookup table 614 indicates that the data source for the requested cell is provided by a synthesizing module 608 (preferably indicated by a pointer to a synthesizing function).

FIG. 7 illustrates exemplary results of table consolidation, triggering, and data synthesis in an embodiment of the present invention. A data table object 700 presents logic table object 710 with a data level table 702 derived from datastore 704. Likewise, data table object 706 presents a logic table object 710 with a data level table 708 derived from

datastore 712. The logic table object 710 consolidates or logically merges all or part of the data level table 702 and data level table 708 into logic level table 714 presented to a logic table object 716. In the illustrated embodiment, the logic level table 714 consists of the logical merger or concatenation of the entire data level table 702 and the entire data level table 708. However, it is likely that the logic table object 710 merely requires a subset of merged table data to satisfy the configuration data request from the caller 718. Likewise, the organization within the rows and columns of the logic level table 714 may differ dramatically from the organization of a mere concatenation of data level tables 702 and 708. The structure of logic level table 714 is determined by the interception/delegation module of the logic table object 710 and a predetermined set of logic component modules, particularly a mapping module.

In the illustrated embodiment, a logic table object 716 overlays the logic table object 710, and presents to the caller 718 with a logic level table 720 that is bigger than the underlying logic level table 714. The two additional columns in table 720 are generated by one or more synthesizing modules in the logic table object 716. Accordingly, the caller 718 can query the column value 722, even though that logic level table's coordinates do not originate from an underlying table or datastore. It should be understood, that in an alternative embodiment, table 720 can have any desired size or configuration in accordance with the logic table object 716.

The logic table object 716 also provides supplemental logic by triggering an external operation in custom activator 724. Custom activators may be related, for example, to activation security (also called launch security) in the COM+ environment that controls which processes can launch a server process. Activation security is automatically applied by the Service Control Manager (SCM) of a particular machine through a custom activator. Upon receipt of a request

from a client to activate an object, the custom activator checks the request against activation-security information stored within the configuration information in the system. Activation security is also checked for activations within the same machine as the client.

FIG. 8 illustrates a flowchart of operations for providing supplemental table logic in an embodiment of the present invention. Operation 800 receives a cursor method call from a caller into the current logic table object. That is, a caller attempts to access data or manipulate the cursor relative to the logic level table presented by the logic table object. Operation 802 determines whether the cursor method call is delegated to an underlying table object. The determination of whether direct delegation (as opposed to interception) that occurs in operation 802 can be accomplished in several ways including (1) receiving the call to the cursor method of the logic table object, which merely calls the corresponding cursor method in one or more underlying table objects; and (2) “fixing up” the vtable associated with the logic table object’s cursor methods to reference an underlying table object’s cursor method directly (i.e., replacing in the vtable the pointer to the logic table object’s cursor method with the pointer to the underlying table object’s cursor method).

In the “fix up” method, the caller calls the cursor method referenced by a pointer in the vtable, anticipating that it is calling the cursor method in the logic table object. Instead, because of the “fix-up”, the cursor method in the underlying table object is executed transparently. After the underlying table object’s method completes, it returns directly to the caller. This optimized method of delegation avoids the cost of the extra indirection in the logic table object.

If delegation does not occur at operation 802, operation 804 intercepts the cursor method call and directs processing to additional logic, which may include without limitation the supplemental logic provided by the module 610, the data generation of the synthesizing

module 608, and the mapping of the mapping module 612 (see FIG. 6). As shown, processing is directed to a supplemental logic module in operation 806, which executes to perform a one or more domain-specific operation (e.g., validation, security enforcement, etc.).

Operation 808 also determines whether delegation occurs. It should be understood that, in one embodiment, operations 802 and 808 are based on the placement of function calls to a cursor method in an underlying table, and that there need be no dynamic determination of delegation at any particular point in the process. Alternatively, however, the delegation can be made conditional on other events or data, wherein operations 802 and 808 may be dynamic determination operations. Furthermore, delegation may occur at any point during the cursor method processing of the logic table object, thereby accommodating supplemental logic pre-processing and/or post-processing for the cursor method of the underlying table object. If there is no delegation at operation 808, then processing returns to the caller in operation 820.

If delegation is determined in operation 802 or 808, operation 810 calls the corresponding cursor method in an underlying cursor method. If the logic table object logically merges more than one lower-level table, then a mapping module is preferably called to determine the appropriate underlying table object or objects that should be called in the delegation.

Operation 812 executes the corresponding cursor method in the underlying table object or objects. Operation 814 returns processing to the current logic table object. Operation 816 determines whether supplemental post-processing is provided. As with delegation operations 802 and 808, operation 816 is preferably provided by the placement of post-processing supplemental logic code in the cursor method of the logic table object, which is executed in operation 818. Operation 820 returns processing to the caller.

FIG. 9 illustrates a flowchart of operations for mapping table data an embodiment of the present invention. The operations relate to a table system such as that illustrated in FIGs. 5 and 6, in which a single logic table object overlays multiple data table objects. Operation 900 receives a cursor method call from a caller requesting access to configuration information.

5 Operation 902 determines the mapping instructions by querying the mapping lookup table to obtain the appropriate underlying table object or table objects and coordinates to which a cursor call should be delegated. Operation 904 calls the corresponding cursor methods in the underlying table objects, providing the mapped coordinates and any other information required by the underlying methods. Operation 906 executes the corresponding cursor methods in the underlying table objects. Operation 908 returns processing to the logic table object.

10 Operation 910 returns processing to the caller.

FIG. 10 illustrates a flowchart of operations for synthesizing data in an embodiment of the present invention. Operation 1000 receives a cursor method call from a caller. The cursor method call specifies functionality relative to a given coordinate in the logic level table. The coordinate may be provided as a parameter in the method call or stored internally in the logic table object. Operation 1002 intercepts the cursor method and references the mapping lookup table to determine the proper mapping for the cursor method call. In the illustrated flowchart, the coordinate does not map to corresponding coordinates in underlying tables, Instead, the mapping module returns a pointer to a synthesizing logic module. Operation 1004 executes the

15 synthesizing logic module to generate data. The data synthesis may rely on data retrieved in operation 1006 from the logic level table, from underlying tables, and from other internal external sources. Operation 1008 performs the synthesis of the data corresponding to the logic

20

level table coordinates associated with the cursor method call. Operation 1010 returns processing to the caller.

The embodiments of the invention described herein are implemented as logical steps in one or more computer systems. The logical operations of the present invention are implemented

5 (1) as a sequence of processor-implemented steps executing in one or more computer systems and (2) as interconnected machine modules within one or more computer systems. The implementation is a matter of choice, dependent on the performance requirements of the computer system implementing the invention. Accordingly, the logical operations making up the embodiments of the invention described herein are referred to variously as operations, steps, 10 objects, or modules.

The above specification, examples and data provide a complete description of the manufacture and use of the composition of the invention. Since many embodiments of the invention can be made without departing from the spirit and scope of the invention, the invention resides in the claims hereinafter appended.

## Claims

### WE CLAIM:

1. A computer program storage medium readable by a computing system and encoding a computer program for executing a computer process providing access to configuration information sourced by at least one datastore, the access being substantially specified by at least one input parameter; the computer process comprising:

5 providing a first level table object instantiated in accordance with a input parameter, the first level table object including a first table-oriented interface having a first table-oriented method;

receiving a call to the first table-oriented method in the first level table object;

10 executing a logic component module responsive to the call, if the first level table object provides domain-specific logic corresponding to the first table-oriented method; and

delegating the call to a corresponding table-oriented method of a lower-level table object to which the first level table object is bound, if the first level table object depends on the lower-level table object to completely service the call.

2. The computer program storage medium of claim 1 wherein the lower-level table object supports a second table-oriented interface identical to the first table-oriented interface of the first level table object.

3. The computer program storage medium of claim 1 wherein the delegating operation comprises:

replacing in a vtable an address to the first table-oriented method in the first level table object with an address to the corresponding table-oriented method of the lower-level table object

4. The computer program storage medium of claim 1 wherein the operation of executing a logic component module comprises:

intercepting the call to the first table-oriented method;

executing the domain-specific logic in a supplemental logic module of the first level table object, responsive to the call received by the first level table object.

5. The computer program storage medium of claim 4 wherein the operation for executing the domain-specific logic comprises:

enforcing complex relationships between a first column and a second column of a logic level table presented to a caller by the first level table object before the datastore is updated.

6. The computer program storage medium of claim 4 wherein the operation for executing the domain-specific logic comprises:

enforcing complex relationships between a logic level table presented to a caller by the first level table object and a second virtual table before the datastore associated with the logic

5 level table is updated.



7. The computer program storage medium of claim 4 wherein the operation for executing the domain-specific logic comprises:

filtering the configuration information accessible by a caller depending on a security level associated with the caller.

8. The computer program storage medium of claim 1 wherein the operation for executing a logic component module comprises:

intercepting the call to the first table-oriented method from a caller, the call being associated with a first coordinate in a logic level table presented to the caller by the first level table object; and

mapping the first coordinate to a second coordinate in a lower-level table presented to the first level table object by the lower-level table object.

9. The computer program storage medium of claim 1 wherein the operation for executing a logic component module comprises:

intercepting the call to the first table-oriented method from a caller, the call being associated with a first coordinate in a logic level table presented to the caller by the table-oriented interface and having no corresponding coordinate in a lower-level table presented by the lower-level table object;

synthesizing data to provide synthesized data associated with the first coordinate in the logic level table; and

returning the synthesized data to the caller.

10. The computer program storage medium of claim 9 wherein the operation for synthesizing data comprises:

accessing lower-level data from at least a second coordinate in the lower-level table; and  
determining the synthesized data based on the lower-level data.

11. The computer program storage medium of claim 1 wherein the operation for executing a logic component module comprises:

triggering an operation external to the first level table object and the lower-level table object.

12. The computer program storage medium of claim 11 wherein the operation for triggering an operation comprises:

triggering a custom activator to provide external activation processing.

13. The computer program storage medium of claim 4 wherein the computer process further comprises:

storing a pointer to the lower-level object usable to access to a lower-level table-oriented method of the lower-level table object.

14. The computer program storage medium of claim 1 wherein the delegating operation comprises:

delegating the call to a corresponding table-oriented method of another lower-level table object to which the first level table object is also bound, if the first level table object depends on

5 the other lower-level table object to completely service the call.

15. The computer program storage medium of claim 1 wherein the computer process further comprises:

    caching read data from the lower-level table object in a read cache of the first level table object.

16. The computer program storage medium of claim 1 wherein the computer process further comprises:

    caching write data intended for the lower-level table object in a write cache of the first level table object.

17. A logic table object, executable by a computer, providing access to configuration information sourced by at least one datastore, the access being substantially specified by at least one input parameter, the logic table object comprising:

    a table-oriented interface including a table-oriented method accessible by a caller to access the configuration information and receiving a call from the caller to the table-oriented method;

    a logic component module providing domain-specific logic to the table-oriented method;

    an interception/delegation module executing the domain-specific logic of the logic component module, responsive to receipt of the call, and further delegating the call to a corresponding table-oriented method of a lower-level table object to which the logic table object is bound, if the logic table object depends on the lower-level table object to completely service the call.

18. The logic table object of claim 17 wherein the logic component module includes a mapping module for translating a first coordinate of a logic level table presented by the logic table object to a second coordinate in a lower-level table presented to the logic table object by the lower-level table object.

19. The logic table object of claim 18 wherein the logic component module includes a mapping lookup table having entries corresponding to coordinates of the logic level table, one or more of the entries including mapping instructions to corresponding coordinates in the lower-level table.

20. The logic table object of claim 17 wherein the logic component module includes a supplemental logic module having domain-specific logic to supplement functionality of the lower-level table object, responsive to the call received by the logic table object

21. The logic table object of claim 20 wherein the supplemental logic module triggers an external operation, responsive to the call.

22. The logic table object of claim 17 wherein the logic component module includes a synthesizing module synthesizing data associated with a first coordinate in a logic level table presented to the caller by the logic table object, wherein no corresponding coordinate exists in a lower-level table presented by the lower-level table object

23. The logic table object of claim 17 wherein the table-oriented interface supported by the logic table object is identical to a second table-oriented interface supported by the lower-level table object to which the logic table object is bound.

24. The logic table object of claim 17 further comprising a first field storing a first pointer to the lower-level object, the pointer being usable to access to a lower-level table-oriented method of the lower-level table object.

25. The logic table object of claim 24 further comprising a second field storing a second pointer to another lower-level table object, and wherein the logic component module comprises a mapping module translating a first coordinate of a logic level table presented by the logic table object to a second coordinate in a lower-level table presented to the logic table object by one of the lower-level table objects.

26. The logic table object of claim 17 further comprising a vtable storing an address to the corresponding table-oriented method of the lower-level table object that corresponds to the table-oriented method of the logic table object called by the caller.

27. The logic table object of claim 17 further comprising:  
a read cache for caching data received from the lower-level table object.

28. The logic table object of claim 17 further comprising:  
a write cache for caching data to be written to the lower-level table object.

29. A computer data signal embodied in a carrier wave by a computing system and encoding a computer program for executing a computer process providing access to requested configuration information through a first level table object including a first table-oriented interface having a first table-oriented method;, the computer program comprising:

- 5 receiving a call to the first table-oriented method in the first level table object;
- intercepting the call to provide supplemental logic, if the first level table object provides domain-specific logic corresponding to the first table-oriented method; and
- delegating the call to a corresponding table-oriented method of a lower-level table object to which the first level table object is bound, if the first level table object depends on the lower-level table object to completely service the call.

30. The computer data signal of claim 29 wherein the first level table object is instantiated in accordance with a input parameter to present a table of the requested configuration data.

31. The computer data signal of claim 29 wherein the computer process further comprises:

- delegating the call to a corresponding table-oriented method of another lower-level table object to which the first level table object is also bound, if the first level table object depends on
- 5 the other lower-level table object to completely service the call.

32. The computer data signal of claim 31 wherein the computer process further comprises:

referencing a mapping lookup table to determine which lower-level table objects are delegated the table-oriented method call.

## **A LOGIC TABLE ABSTRACTION LAYER FOR ACCESSING CONFIGURATION INFORMATION**

### **Abstract of the Invention**

A logic table object for accessing configuration information sourced by one or more  
5 datastores is employed in a catalog environment of an attribute-based programming model. A  
logic table object may provide supplemental logic, consolidate multiple underlying tables, map  
between different disparate table coordinates, trigger external operations, and synthesize data for  
inclusion in a virtual table presented to a caller. A logic table object can provide access to  
configuration information cached in an underlying data table object and derived from a datastore.  
10 Alternatively, a logic table can include its own cache, servicing calls therefrom until an update to  
or from an underlying datastore is required.

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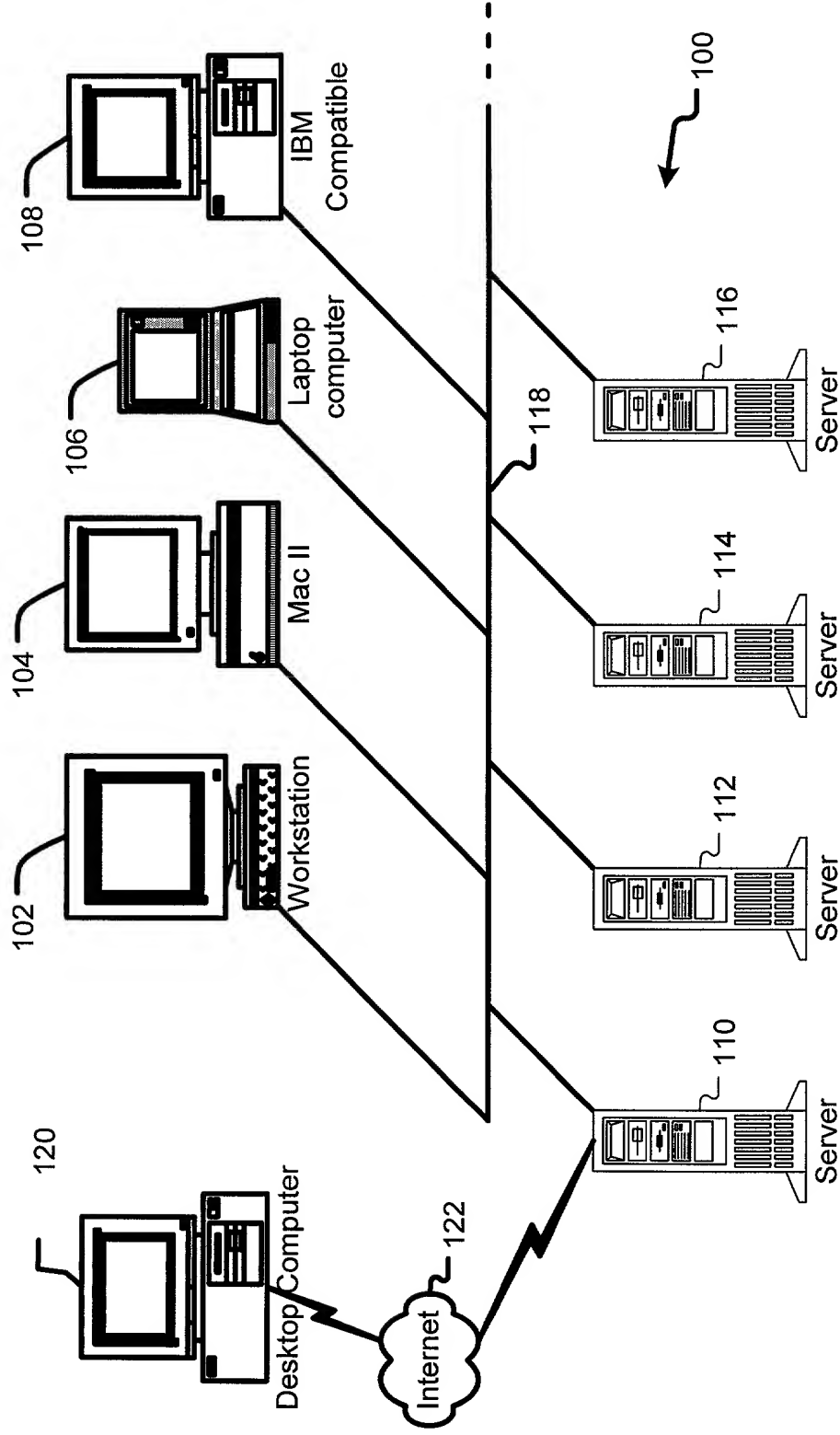


FIG. 1

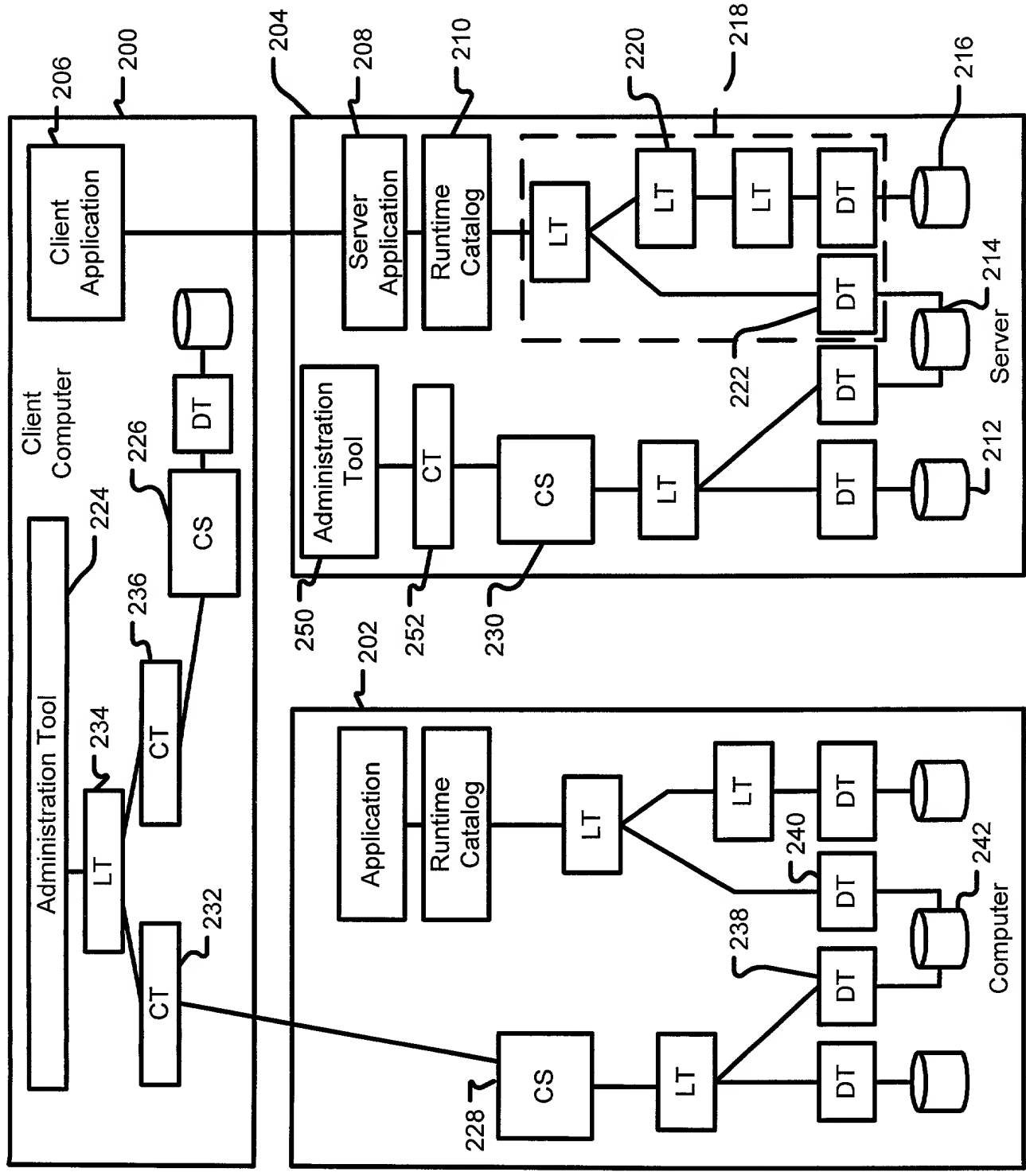


FIG. 2

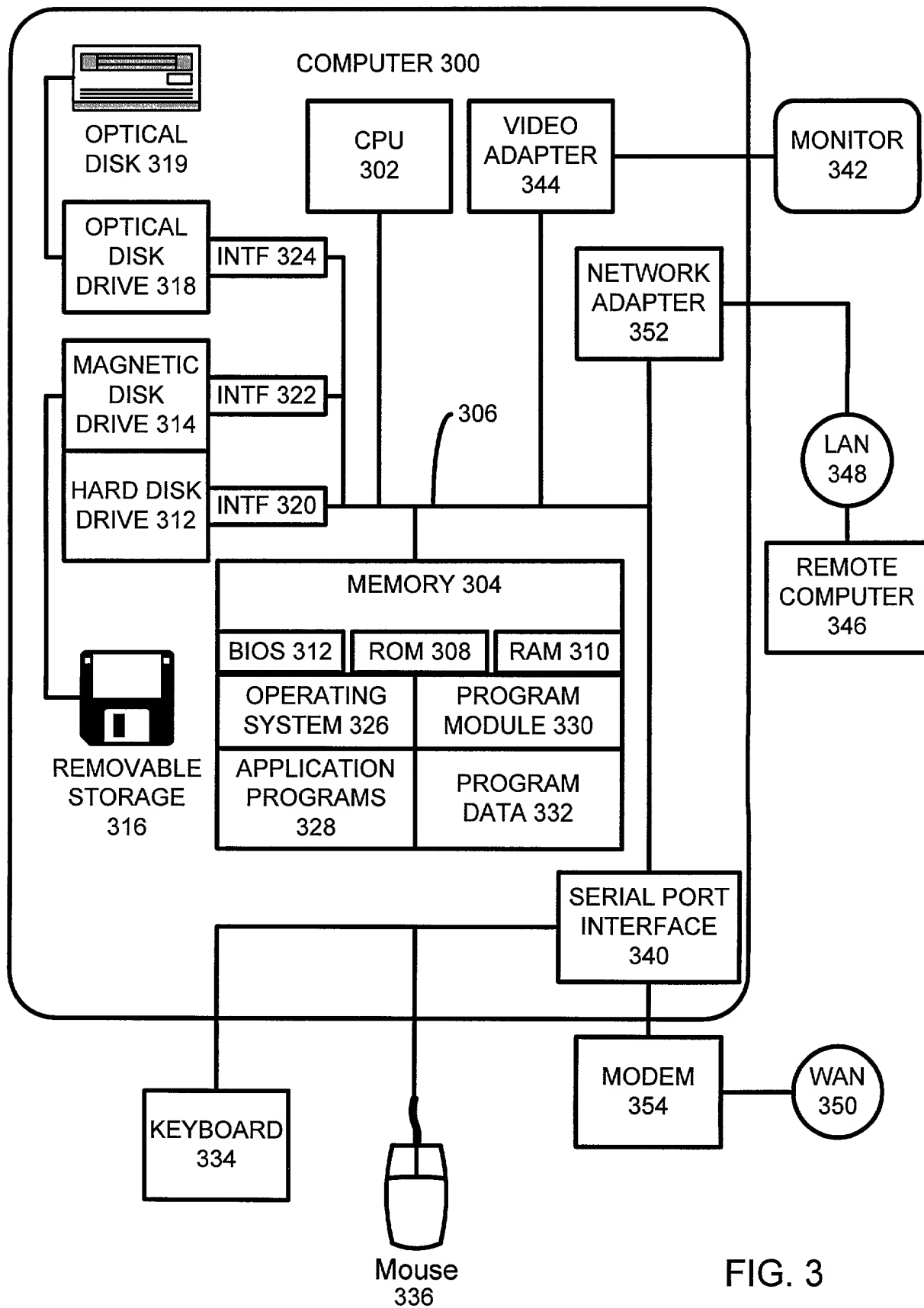


FIG. 3

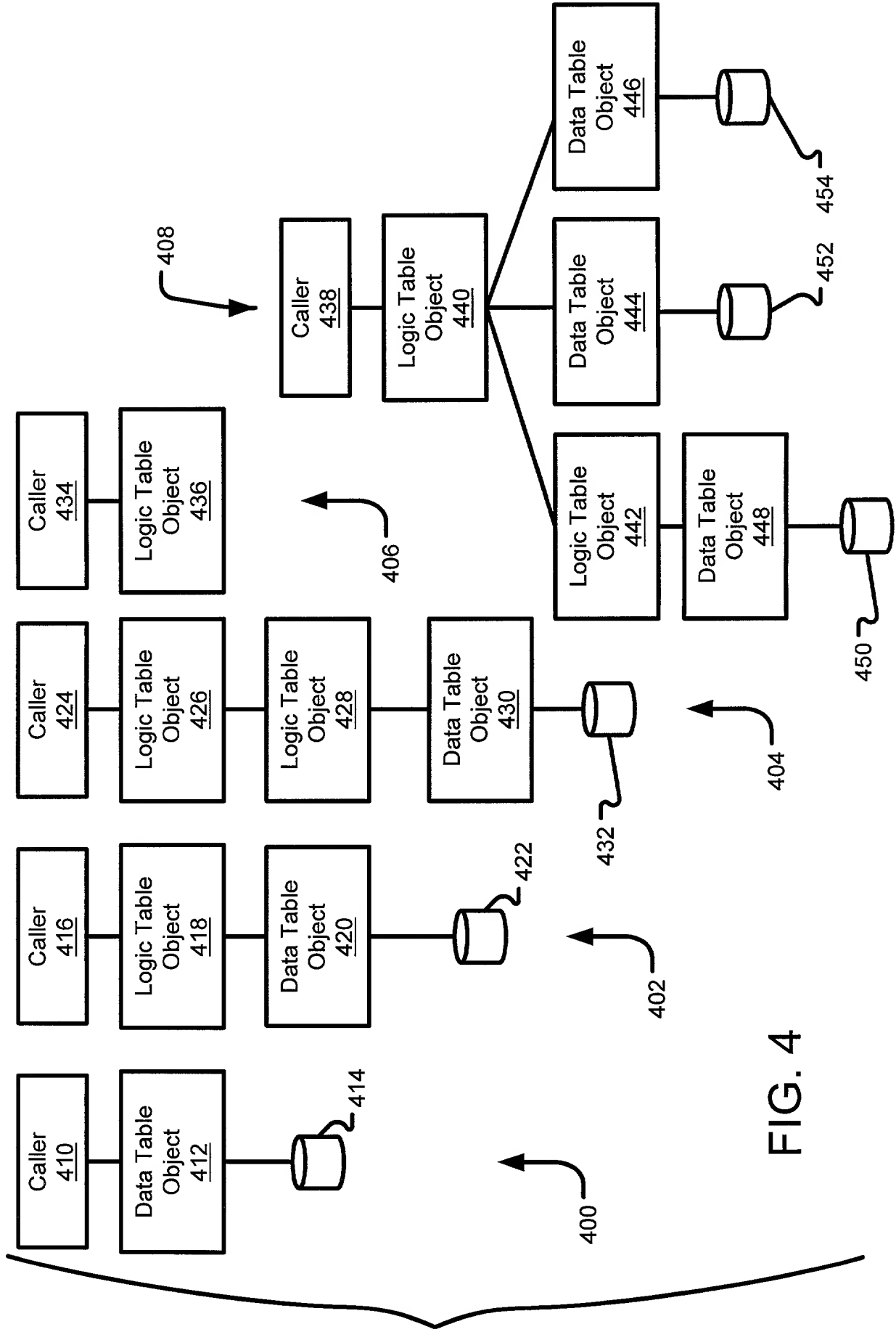


FIG. 4

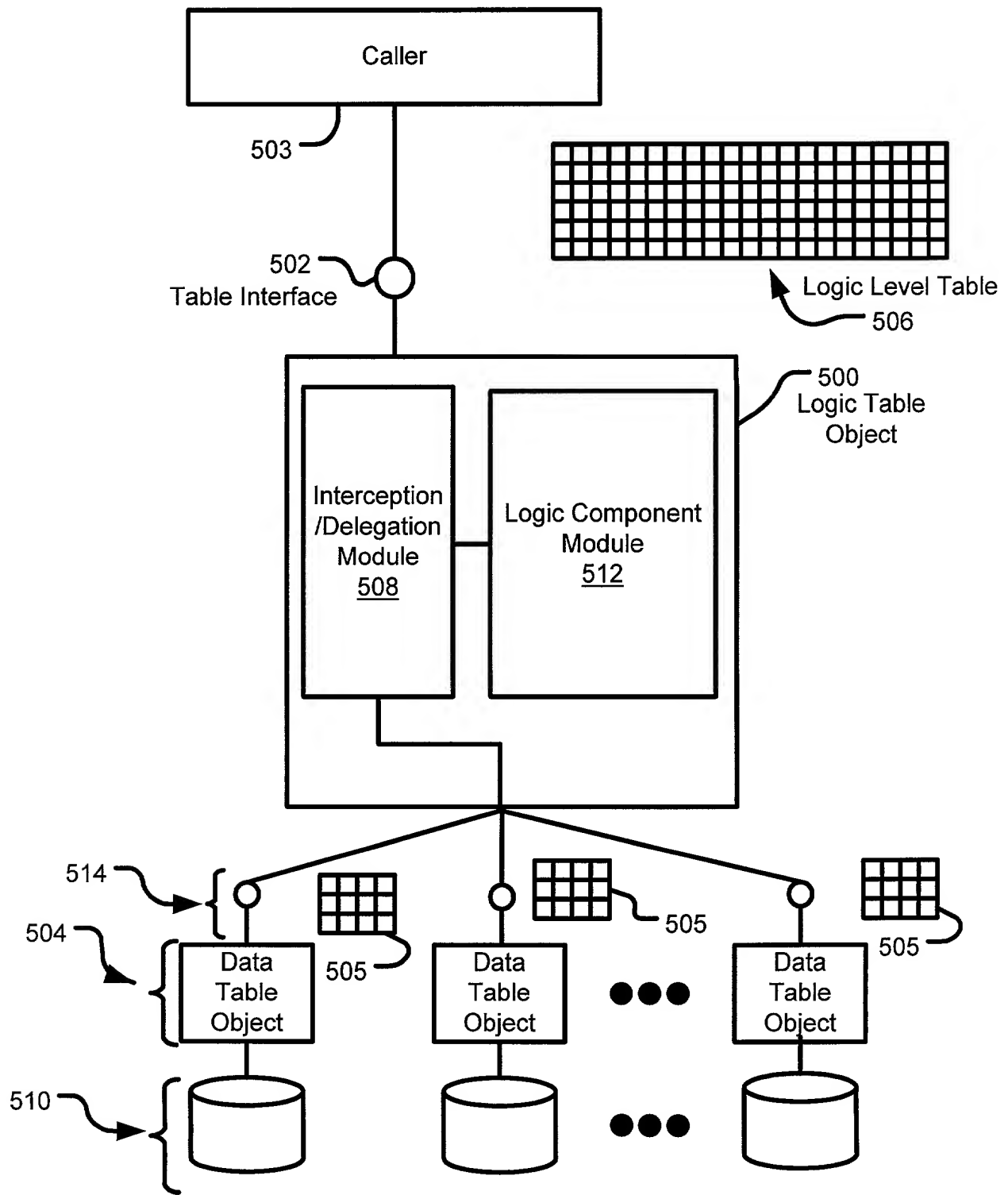


FIG. 5

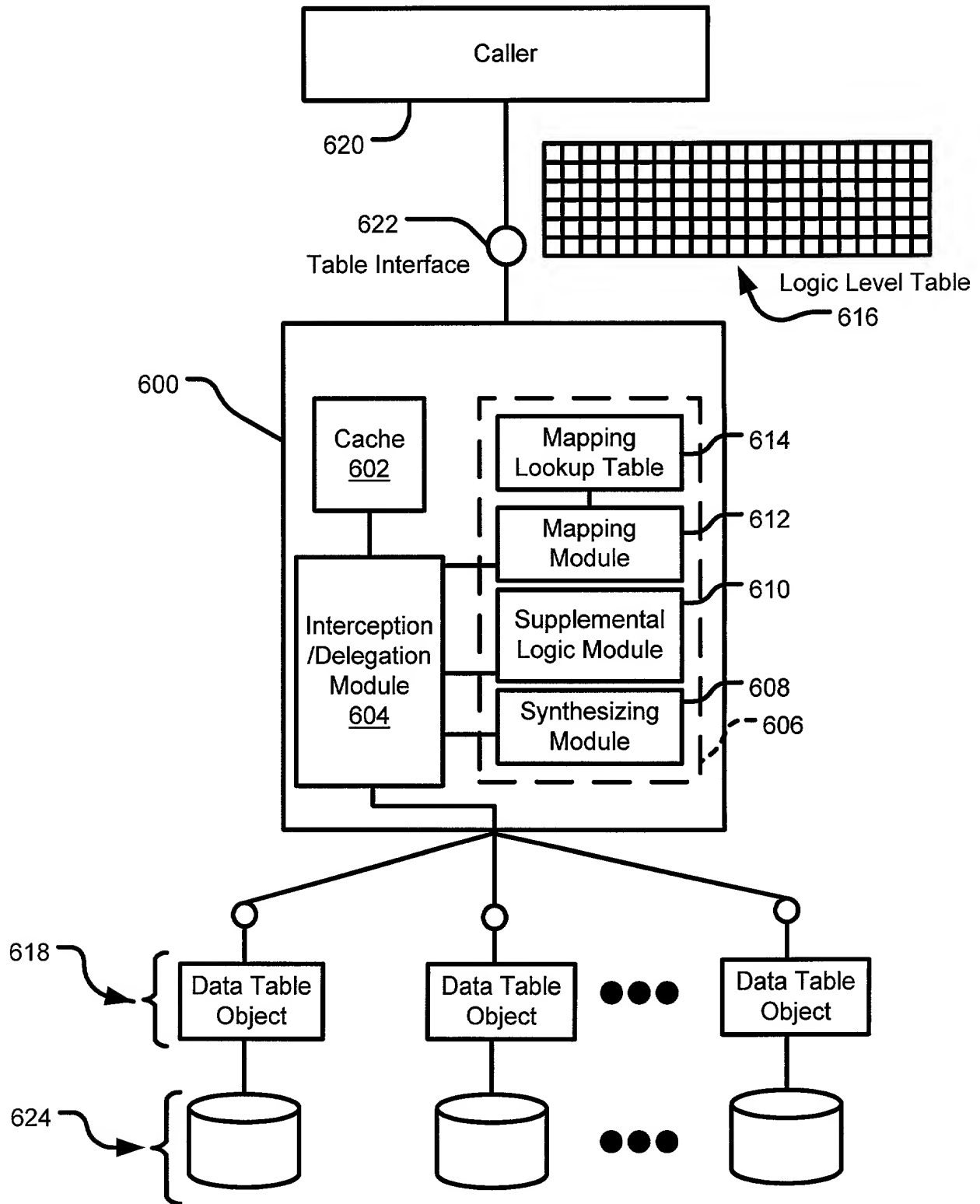


FIG. 6

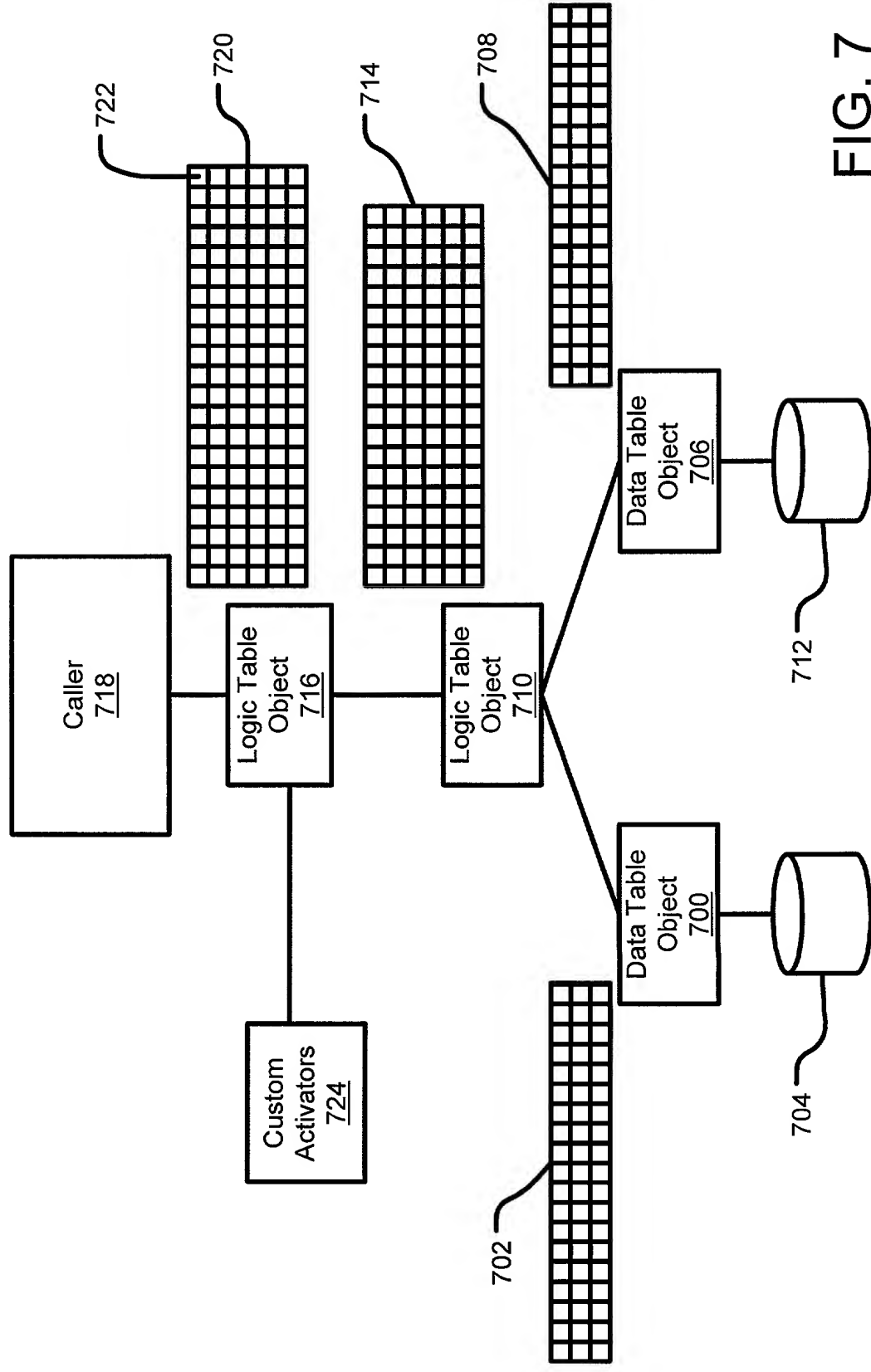


FIG. 7

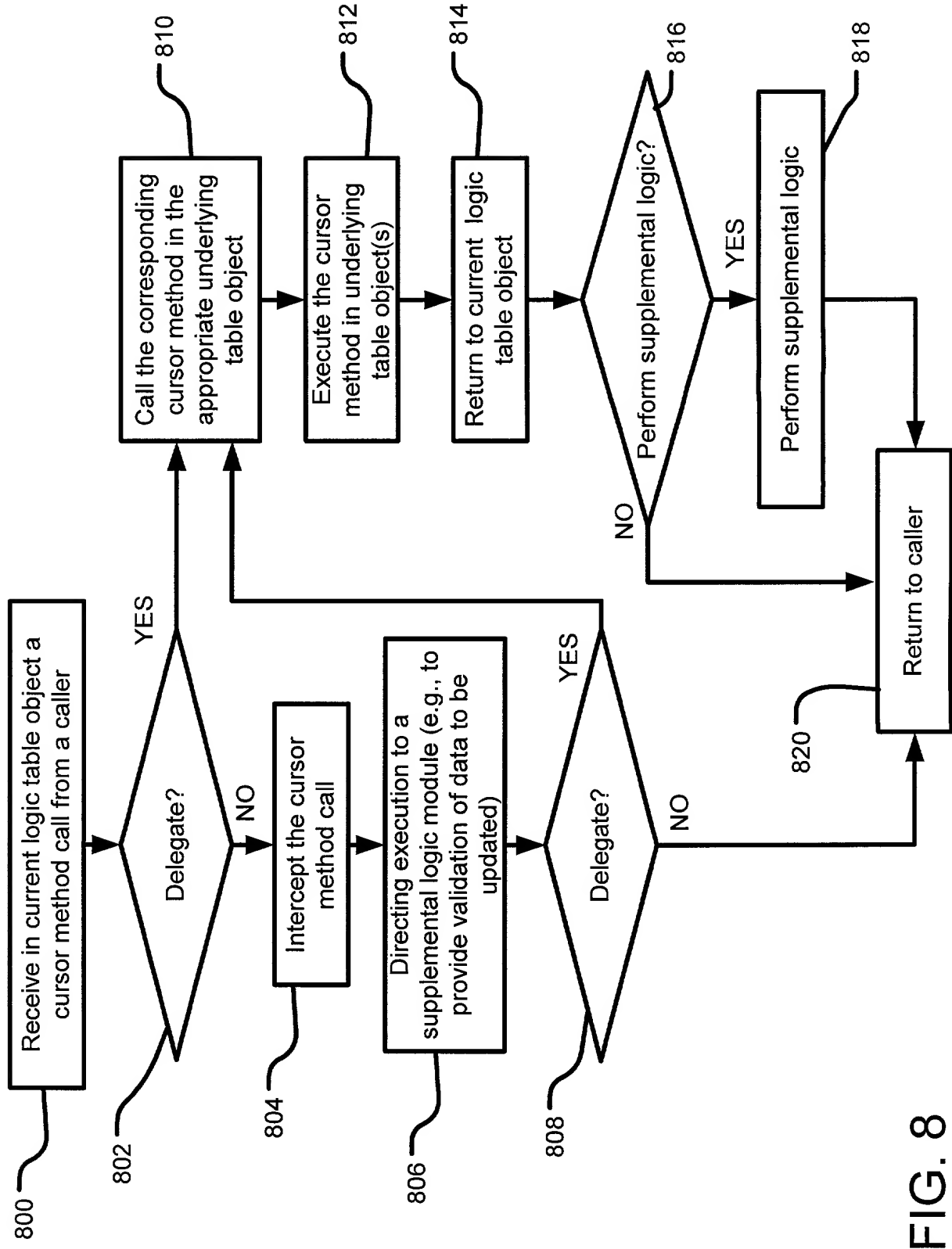


FIG. 8



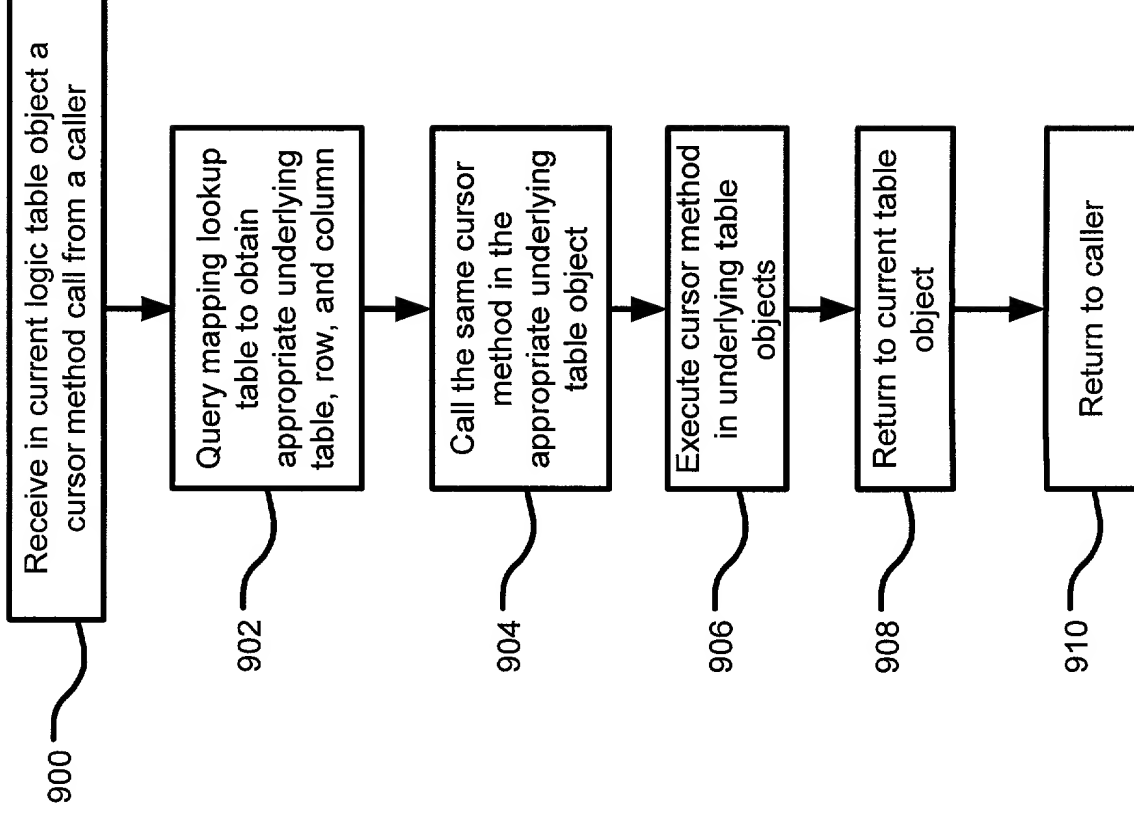


FIG. 9

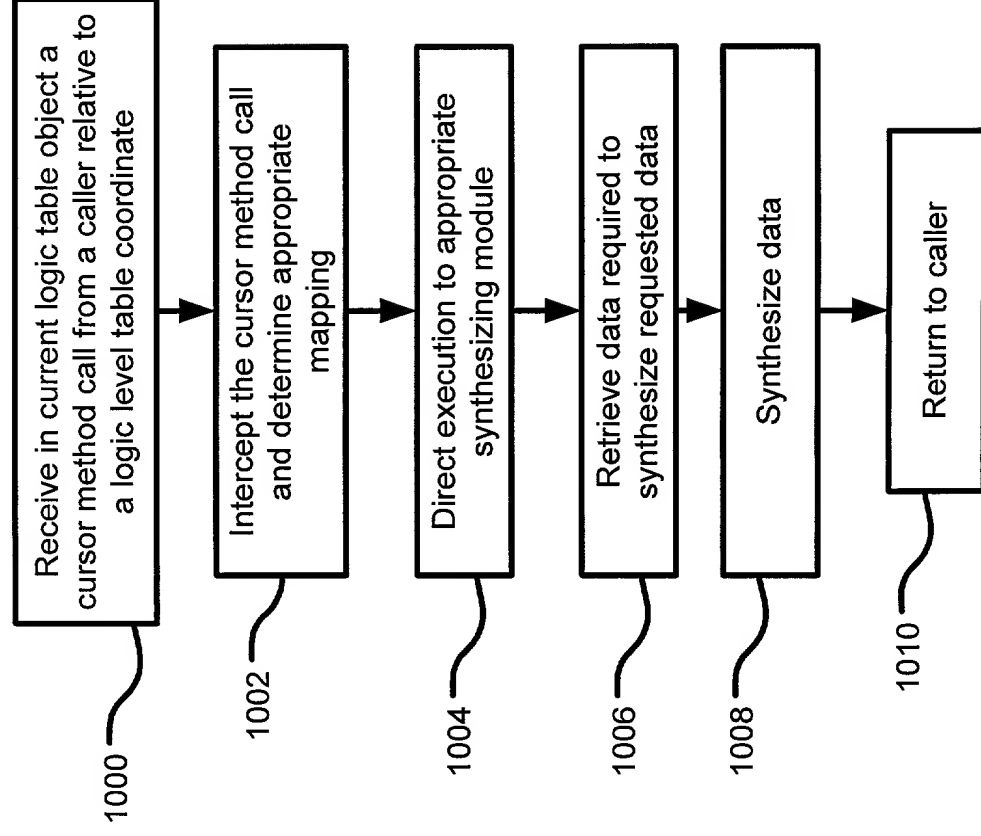


FIG. 10

MERCHANT & GOULD P.C.

**United States Patent Application**

**COMBINED DECLARATION AND POWER OF ATTORNEY**

As a below named inventor I hereby declare that my residence, post office address and citizenship are as stated below next to my name; that

I verily believe I am the original, first and a sole inventor of the subject matter which is claimed and for which a patent is sought on the invention entitled A LOGIC TABLE ABSTRACTION LAYER FOR ACCESSING CONFIGURATION INFORMATION, the specification of which I have reviewed and for which I solicit a United States patent.

I hereby state that I have reviewed and understand the contents of the above-identified specification, including the claims.

I acknowledge the duty to disclose information which is material to the patentability of this application in accordance with Title 37, Code of Federal Regulations, § 1.56 (attached hereto).

I hereby claim foreign priority benefits under Title 35, United States Code, § 119/365 of any foreign application(s) for patent or inventor's certificate listed below and have also identified below any foreign application for patent or inventor's certificate having a filing date before that of the application on the basis of which priority is claimed:

a. ☒ no such applications have been filed.

b. ☐ such applications have been filed as follows:

FOREIGN APPLICATION(S), IF ANY, CLAIMING PRIORITY UNDER 35 USC § 119			
COUNTRY	APPLICATION NUMBER	DATE OF FILING (day, month, year)	DATE OF ISSUE (day, month, year)
ALL FOREIGN APPLICATION(S), IF ANY, FILED BEFORE THE PRIORITY APPLICATION(S)			
COUNTRY	APPLICATION NUMBER	DATE OF FILING (day, month, year)	DATE OF ISSUE (day, month, year)

I hereby claim the benefit under Title 35, United States Code, § 120/365 of any United States and PCT international application(s) listed below and, insofar as the subject matter of each of the claims of this application is not disclosed in the prior United States application in the manner provided by the first paragraph of Title 35, United States Code, § 112, I acknowledge the duty to disclose material information as defined in Title 37, Code of Federal Regulations, § 1.56(a) which occurred between the filing date of the prior application and the national or PCT international filing date of this application.

U.S. APPLICATION NUMBER	DATE OF FILING (day, month, year)	STATUS (patented, pending, abandoned)

I hereby claim the benefit under Title 35, United States Code § 119(e) of any United States provisional application(s) listed below:

U.S. PROVISIONAL APPLICATION NUMBER	DATE OF FILING (Day, Month, Year)

I hereby appoint the following attorney(s) and/or patent agent(s) to prosecute this application and to transact all business in the Patent and Trademark Office connected herewith:

Albrecht, John W.	Reg. No. 40,481	Lacy, Paul E.	Reg. No. 38,946
Anderson, Gregg I.	Reg. No. 28,828	Larson, James A.	Reg. No. 40,443
Ansems, Gregory M.	Reg. No. 42,264	Lasky, Michael B.	Reg. No. 29,555
Batzli, Brian H.	Reg. No. 32,960	Liepa, Mara E.	Reg. No. 40,066
Beard, John L.	Reg. No. 27,612	Lindquist, Timothy A.	Reg. No. 40,701
Black, Bruce E.	Reg. No. 41,622	McDonald, Daniel W.	Reg. No. 32,044
Bruess, Steven C.	Reg. No. 34,130	McIntyre, Iain A.	Reg. No. 40,337
Byrne, Linda M.	Reg. No. 32,404	Mueller, Douglas P.	Reg. No. 30,300
Carlson, Alan G.	Reg. No. 25,959	Nelson, Albin J.	Reg. No. 28,650
Caspers, Philip P.	Reg. No. 33,227	Pauly, Daniel M.	Reg. No. 40,123
Chiapetta, James R.	Reg. No. 39,634	Phillips, John B.	Reg. No. 37,206
Clifford, John A.	Reg. No. 30,247	Plunkett, Theodore	Reg. No. 37,209
Cochran, William W.	Reg. No. 26,652	Pytel, Melissa J.	Reg. No. 41,512
Daignault, Ronald A.	Reg. No. 25,968	Reich, John C.	Reg. No. 37,703
Daley, Dennis R.	Reg. No. 34,994	Reiland, Earl D.	Reg. No. 25,767
Dalglish, Leslie E.	Reg. No. 40,579	Schmaltz, David G.	Reg. No. 39,828
Daulton, Julie R.	Reg. No. 36,414	Schuman, Mark D.	Reg. No. 31,197
Devries Smith, Katherine M.	Reg. No. 42,157	Schumann, Michael D.	Reg. No. 30,422
DiPietro, Mark J.	Reg. No. 28,707	Scull, Timothy B.	Reg. No. 42,137
Edell, Robert T.	Reg. No. 20,187	Sebald, Gregory A.	Reg. No. 33,280
Epp Ryan, Sandra	Reg. No. 39,667	Skoog, Mark T.	Reg. No. 40,178
Funk, Steven R.	Reg. No. 37,830	Soderberg, Richard	Reg. No. P43,352
Glance, Robert J.	Reg. No. 40,620	Storer, Shelley D.	Reg. No. P45,135
Golla, Charles E.	Reg. No. 26,896	Sumner, John P.	Reg. No. 29,114
Gorman, Alan G.	Reg. No. 38,472	Sumners, John S.	Reg. No. 24,216
Gould, John D.	Reg. No. 18,223	Swenson, Erik G.	Reg. No. P45,147
Gregson, Richard	Reg. No. 41,804	Tellekson, David K.	Reg. No. 32,314
Gresens, John J.	Reg. No. 33,112	Trembath, Jon R.	Reg. No. 38,344
Hamre, Curtis B.	Reg. No. 29,165	Underhill, Albert L.	Reg. No. 27,403
Hillson, Randall A.	Reg. No. 31,838	Vandenburgh, J. Derek	Reg. No. 32,179
Holzer, Jr., Richard J.	Reg. No. 42,668	Wahl, John R.	Reg. No. 33,044
Johnston, Scott W.	Reg. No. 39,721	Welter, Paul A.	Reg. No. 20,890
Kadjevitch, Natalie D.	Reg. No. 34,196	Whipps, Brian	Reg. No. 43,261
Kastelic, Joseph M.	Reg. No. 37,160	Wickhem, J. Scot	Reg. No. 41,376
Kettelberger, Denise	Reg. No. 33,924	Williams, Douglas J.	Reg. No. 27,054
Knearl, Homer L.	Reg. No. 21,197	Witt, Jonelle	Reg. No. 41,980
Kowalchuk, Alan W.	Reg. No. 31,535	Xu, Min S.	Reg. No. 39,536
Kowalchuk, Katherine M.	Reg. No. 36,848		

In addition, I also hereby appoint the following attorneys to prosecute this application and to transact all business in the U.S. Patent and Trademark Office in connection therewith:

Katie E. Sako, Reg. No. 32,628  
Daniel D. Crouse, Reg. No. 32,022

I hereby authorize them to act and rely on instructions from and communicate directly with the person/assignee/attorney/firm/ organization who/which first sends/sent this case to them and by whom/which I hereby declare that I have consented after full disclosure to be represented unless/until I instruct Merchant & Gould P.C. to the contrary.

Please direct all correspondence in this case to Merchant & Gould P.C. at the address indicated below:

Homer L. Knearl  
Merchant & Gould P.C.  
3100 Norwest Center  
90 South Seventh Street  
Minneapolis, MN 55402-4131  
303.357.1633

I hereby declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under Section 1001 of Title 18 of the United States Code and that such willful false statements may jeopardize the validity of the application or any patent issued thereon.

201	<b>Full Name Of Inventor</b>	<b>Family Name</b> Craig	<b>First Given Name</b> Robert	<b>Second Given Name</b> M.
	<b>Residence &amp; Citizenship</b>	<b>City</b> Bellevue	<b>State or Foreign Country</b> WA	<b>Country of Citizenship</b> USA
	<b>Post Office Address</b>	<b>Post Office Address</b> 2510 165 <sup>th</sup> Avenue NE	<b>City</b> Bellevue	<b>State &amp; Zip Code/Country</b> WA 98008/USA
<b>Signature of Inventor 201:</b>			<b>Date:</b>	

### § 1.56 Duty to disclose information material to patentability.

(a) A patent by its very nature is affected with a public interest. The public interest is best served, and the most effective patent examination occurs when, at the time an application is being examined, the Office is aware of and evaluates the teachings of all information material to patentability. Each individual associated with the filing and prosecution of a patent application has a duty of candor and good faith in dealing with the Office, which includes a duty to disclose to the Office all information known to that individual to be material to patentability as defined in this section. The duty to disclose information exists with respect to each pending claim until the claim is canceled or withdrawn from consideration, or the application becomes abandoned. Information material to the patentability of a claim that is canceled or withdrawn from consideration need not be submitted if the information is not material to the patentability of any claim remaining under consideration in the application. There is no duty to submit information which is not material to the patentability of any existing claim. The duty to disclose all information known to be material to patentability is deemed to be satisfied if all information known to be material to patentability of any claim issued in a patent was cited by the Office or submitted to the Office in the manner prescribed by §§ 1.97(b)–(d) and 1.98. However, no patent will be granted on an application in connection with which fraud on the Office was practiced or attempted or the duty of disclosure was violated through bad faith or intentional misconduct. The Office encourages applicants to carefully examine:

- (1) prior art cited in search reports of a foreign patent office in a counterpart application, and
- (2) the closest information over which individuals associated with the filing or prosecution of a patent application believe any pending claim patentably defines, to make sure that any material information contained therein is disclosed to the Office.

(b) Under this section, information is material to patentability when it is not cumulative to information already of record or being made of record in the application, and

- (1) It establishes, by itself or in combination with other information, a prima facie case of unpatentability of a claim;
- or
- (2) It refutes, or is inconsistent with, a position the applicant takes in:
    - (i) Opposing an argument of unpatentability relied on by the Office, or
    - (ii) Asserting an argument of patentability.

A prima facie case of unpatentability is established when the information compels a conclusion that a claim is unpatentable under the preponderance of evidence, burden-of-proof standard, giving each term in the claim its broadest reasonable construction consistent with the specification, and before any consideration is given to evidence which may be submitted in an attempt to establish a contrary conclusion of patentability.

(c) Individuals associated with the filing or prosecution of a patent application within the meaning of this section are:

- (1) Each inventor named in the application:

(2) Each attorney or agent who prepares or prosecutes the application; and

(3) Every other person who is substantively involved in the preparation or prosecution of the application and who is associated with the inventor, with the assignee or with anyone to whom there is an obligation to assign the application.

(d) Individuals other than the attorney, agent or inventor may comply with this section by disclosing information to the attorney, agent, or inventor.